

A GENTLE INTRODUCTION TO EPISTEMIC PLANNING

FOUNDATIONS AND CHALLENGES

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A Picture from the Top

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Epistemic Logic

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Dynamic Epistemic Logic

A (SLIGHTLY) PHILOSOPHICAL INTRODUCTION

To Know or to Believe? ...

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The Tripartite Analysis of Knowledge

S **knows** that p iff

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⇒ **Justified True Belief (JTB)**

Is Justified True Belief actually Knowledge?

John is standing outside a field and, within it, he sees what looks exactly like a sheep.

→ Does John **know** that there is a sheep in the field?

Let's analyse the situation:

- 1 John sure **believes** that there is a sheep in the field.
- 2 John is also justified in believing so: he clearly sees it!
- 3 But **is it true** that there is a sheep in the field?

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Moreover, there is actually a sheep behind the hill in the middle of the field.

→ What can we say now?

EPISTEMIC LOGIC

Let \mathcal{P} be a finite set of **propositional atoms** and $\mathcal{AG} = \{1, \dots, n\}$ a finite set of **agents**. The **language** $\mathcal{L}_{\mathcal{P}, \mathcal{AG}}$ of **Epistemic Logic** is given by the BNF:

Definition (Language of Epistemic Logic)

$$\varphi ::= p \mid \neg\varphi \mid \varphi \wedge \varphi \mid \Box_i\varphi,$$

- Operator \Box_i : depending on the context, describes what agent i **knows** or **believes**.
- Dual operator \Diamond_i : describes what agent i **considers to be possible** or **compatible**.

An *epistemic state* represents both **factual** information and what agents **know/believe**.

Definition (Epistemic Model)

An *epistemic model* is a triple $M = (W, R, V)$, where:

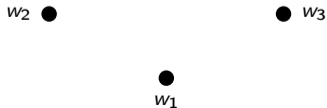
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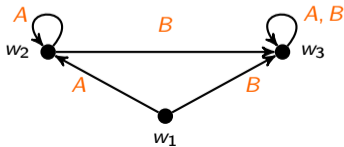
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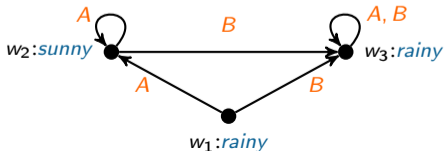
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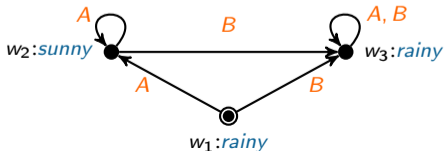
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An *epistemic state* is a pair (M, W_d) s.t. $W_d \subseteq W$ is a non-empty set of **designated worlds**.

Example



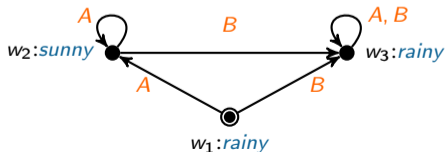
Definition (Truth)

Let $s = (M, W_d)$, where $M = (W, R, V)$, be an *epistemic state* and let $w \in W$:

$(M, w) \models p$	iff	$w \in V(p)$
$(M, w) \models \neg\varphi$	iff	$(M, w) \not\models \varphi$
$(M, w) \models \varphi \wedge \psi$	iff	$(M, w) \models \varphi$ and $(M, w) \models \psi$
$(M, w) \models \Box_i\varphi$	iff	$\forall v$ if wR_iv then $(M, v) \models \varphi$

Moreover, $(M, W_d) \models \varphi$ iff $\forall w$ if $w \in W_d$ then $(M, w) \models \varphi$.

Example



- $\Box_{Anne} \textit{sunny}$
- $\Box_{Bob} \textit{rainy}$
- $\Box_{Anne} \Box_{Bob} \textit{rainy}$
- $\Diamond_{Bob} \Box_{Anne} \textit{rainy}$

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	Axiom	Frame Property	Knowledge	Belief
K	$\Box_i(\varphi \rightarrow \psi) \rightarrow (\Box_i\varphi \rightarrow \Box_i\psi)$	-	✓	✓
T	$\Box_i\varphi \rightarrow \varphi$	Reflexivity	✓	
D	$\Box_i\varphi \rightarrow \Diamond_i\varphi$	Seriality	✓	✓
4	$\Box_i\varphi \rightarrow \Box_i\Box_i\varphi$	Transitivity	✓	✓
5	$\neg\Box_i\varphi \rightarrow \Box_i\neg\Box_i\varphi$	Euclideaness	✓	✓

An **epistemic state** represents:

- **Knowledge**, when it satisfies axioms **K**, **T**, **4** and **5** \Rightarrow **Logic S5_n**
- **Belief**, when it satisfies axioms **K**, **D**, **4** and **5** \Rightarrow **Logic KD45_n**

DYNAMIC EPISTEMIC LOGIC

Actions in Classical Planning

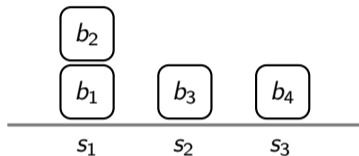
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- 1 Propositional
- 2 Single-agent
- 3 Fully Observable
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Example (Blocks World)

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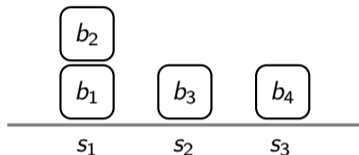
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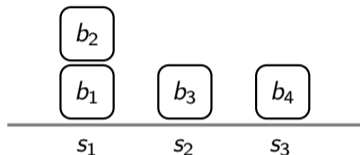
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Action $move(b, x, y)$:

- $Pre(move(b, x, y)) = On(b, x) \wedge Clear(b) \wedge Clear(y)$
- $Eff(move(b, x, y)) = \{On(b, y), Clear(x), \neg On(b, x), \neg Clear(y)\} \triangleright \top$

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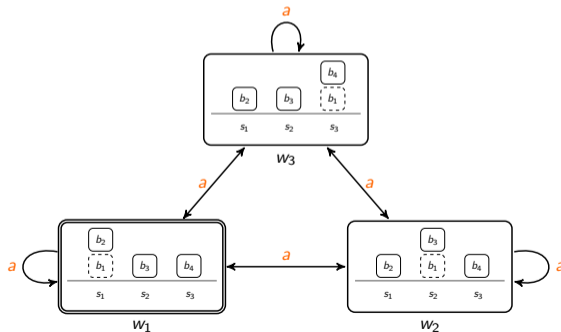
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→ We now incrementally move from classical actions to epistemic actions.

Epistemic Blocks World

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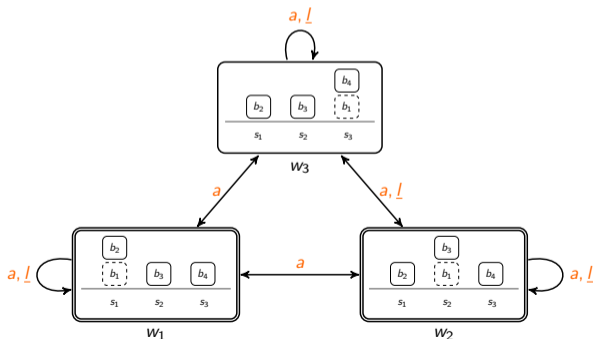
- Agent a : only sees from above.



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Example (Multi-Agent Epistemic Blocks World)

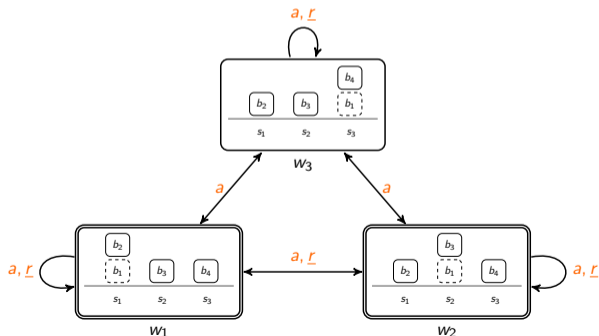
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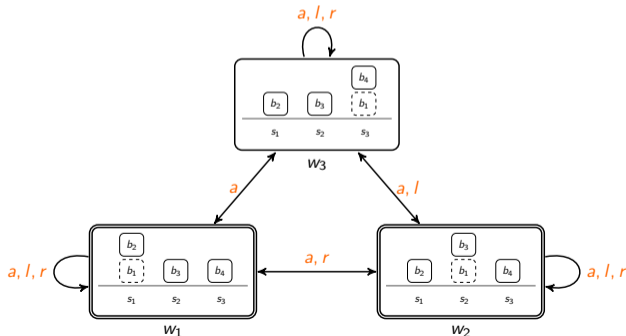
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An *event model* is a quadruple $\mathcal{E} = (E, Q, pre, post)$, where:

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Intuitively:

- An **event** can be seen as a **classical action**.
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Given (M, W_d) and (\mathcal{E}, E_d) , where $M = (W, R, V)$ and $\mathcal{E} = (E, Q, pre, post)$, their **product update** $(M, W_d) \otimes (\mathcal{E}, E_d)$ is the **epistemic state** $((W', R', V'), W'_d)$ where:

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Public Announcement

Agent r tells everybody that he knows that $\neg On(b_1, s_3)$.



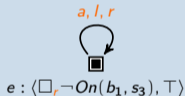
The diagram shows a small square representing an agent. Above it, the text "a, l, r" is written in orange. A curved arrow starts from the top of the square and points back to the top, indicating a self-announcement or a public announcement to all agents.

$$e : \langle \Box_r \neg On(b_1, s_3), \top \rangle$$

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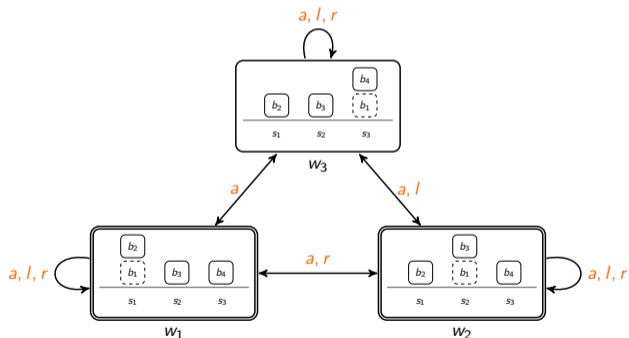
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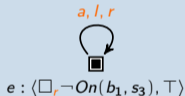
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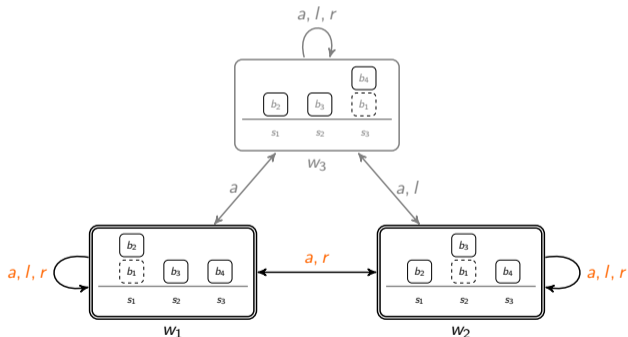
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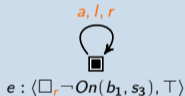
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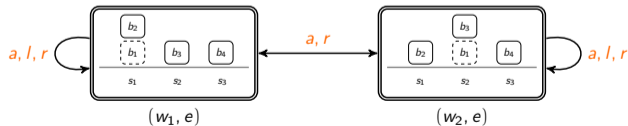
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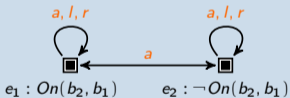
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- Agents r and l observe what is actually being sensed.
- Agent a can not directly observe what agent r is seeing.



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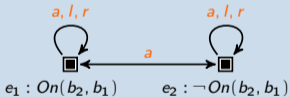
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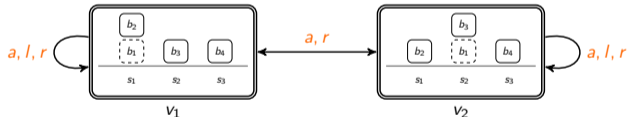
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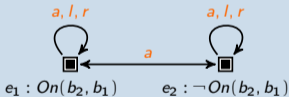
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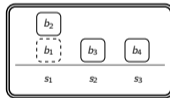
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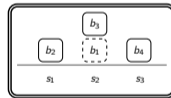
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Trivial postconditions are omitted.



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(v_2, e_2)

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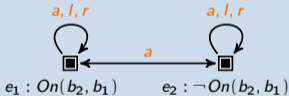
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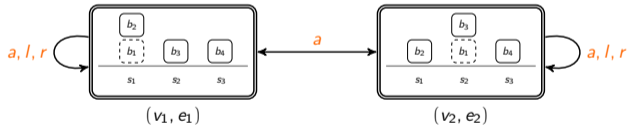
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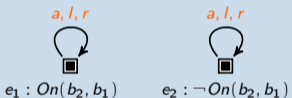
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→ **Non-deterministic** action!



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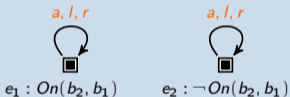
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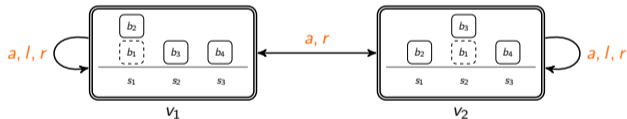
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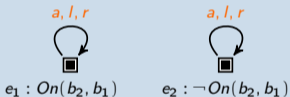
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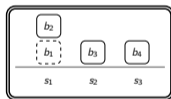
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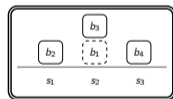
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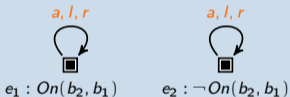
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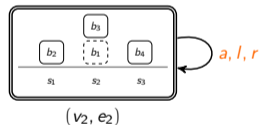
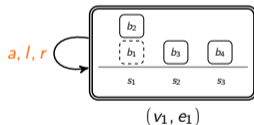
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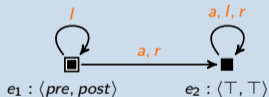
Private Ontic Actions

Example

Private Ontic Action

Agent *l* privately moves block b_2 from b_1 to b_3 , where:

- $pre = On(b_2, b_1) \wedge Clear(b_2) \wedge Clear(b_3)$
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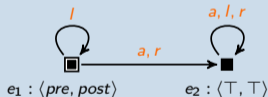
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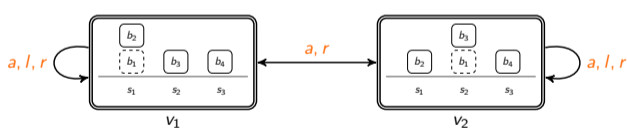
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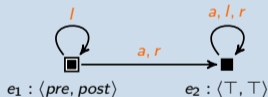
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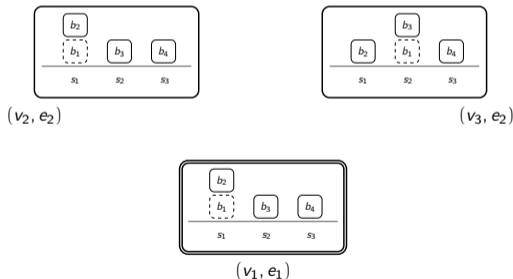
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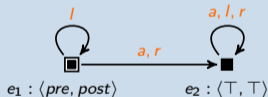


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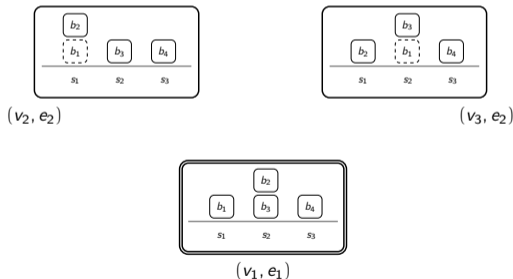
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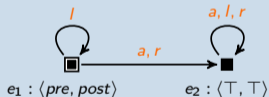
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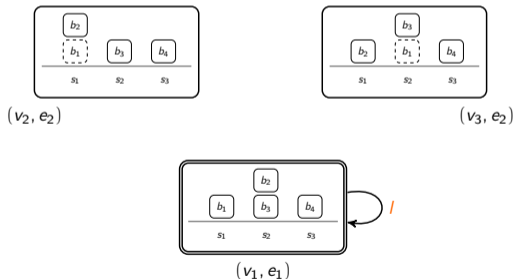
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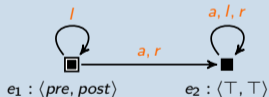


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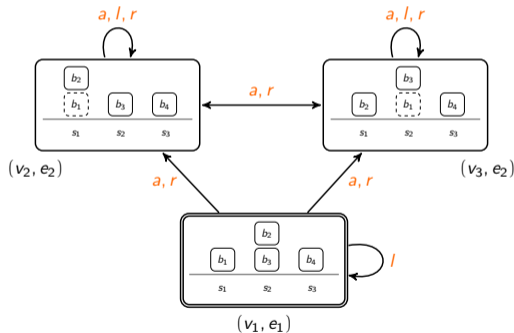
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Classical Vs. Epistemic Actions

To summarize:

Classical actions are:

- 1 Propositional
- 2 Single-agent
- 3 Fully Observable
- 4 Deterministic

Epistemic actions are:

- 1 Modal
- 2 Multi-agent
- 3 Partially Observable
- 4 Non-deterministic

Moreover, epistemic actions model both factual and higher-order knowledge change.

→ There are **no restrictions** on the reasoning power of agents! (More on this later)

Definition (Planning Task)

An **(epistemic) planning task** is a triple $T = (s_0, \mathcal{A}, \varphi_g)$, where:

- s_0 is an **initial epistemic state**;
- \mathcal{A} is a finite **set of actions**;
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Definition (Solution)

A **solution** to a planning task $(s_0, \mathcal{A}, \varphi_g)$ is a finite sequence $\alpha_1, \dots, \alpha_m$ of actions of \mathcal{A} s.t.:

- 1 For each $1 \leq k \leq m$, α_k is applicable in $s_0 \otimes \alpha_1 \otimes \dots \otimes \alpha_{k-1}$, and
- 2 $s_0 \otimes \alpha_1 \otimes \dots \otimes \alpha_m \models \varphi_g$.

Definition (Plan Existence Problem)

Let $n \geq 1$ and \mathcal{T} be a class of planning tasks. $\text{PlanEx}(\mathcal{T}, n)$ is the following decision problem: “Given a planning task $T = (s_0, \mathcal{A}, \varphi_g) \in \mathcal{T}$, where $|\mathcal{AG}| = n$, does T have a solution?”

Theorem (Bolander and Andersen [BA11])

*Let \mathcal{T} be the class of all epistemic planning tasks and let $n \geq 1$. Then, $\text{PlanEx}(\mathcal{T}, n)$ is **undecidable**.*

CURRENT CHALLENGES

Decidable Fragments

A great deal of effort has been spent over the past decade to devise **decidable** fragments of the epistemic plan existence problem.

Let $\mathcal{T}(a, b)$ denote the class of epistemic planning tasks where:

- a is the maximum modal depth of **preconditions**, and
- b is the maximum modal depth of **postconditions**. We indicate with $b = -1$ the absence of postconditions.

$\text{PlanEx}(\mathcal{T}(0, -1), n)$	PSPACE-complete [CMS16]
$\text{PlanEx}(\mathcal{T}(1, -1), n)$	Unknown [CMS16]
$\text{PlanEx}(\mathcal{T}(2, -1), n)$	UNDECIDABLE [CMS16]
$\text{PlanEx}(\mathcal{T}(0, 0), n)$	DECIDABLE [YWL13; AMP14]
$\text{PlanEx}(\mathcal{T}(1, 0), n)$	DECIDABLE [Bol+20]

Decidable Fragments (cont.)

Others have focused on considering the plan existence problem of tasks under well-known modal logics (Aucher and Bolander [AB13]).

Logic	Single-agent	Multi-agent
K	UNDECIDABLE	UNDECIDABLE
KT		
K4		
K45	DECIDABLE	
S4	UNDECIDABLE	
S5	DECIDABLE	

A Semantic Approach

What if we combined the two previous approaches together?

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$$\mathbf{C} \quad \Box_i \Box_j \varphi \rightarrow \Box_j \Box_i \varphi$$

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Lemma (Burigana et al. [Bur+23])

Let (M, W_d) be a bisimulation-contracted C-S5_n-state, with $M = (W, R, V)$. Then, $|W|$ is **bounded in n and $|\mathcal{P}|$** .

Theorem (Burigana et al. [Bur+23])

The plan existence problem in **C-S5_n** is **decidable**.

Generalizing Commutativity

Let $b > 1$ be a fixed integer constant:

b -Commutativity

$$\mathbf{C}^b \quad (\Box_i \Box_j)^b \varphi \rightarrow (\Box_j \Box_i)^b \varphi$$

Let $1 < \ell \leq n$ be a fixed integer constant, let $\langle i_1, \dots, i_\ell \rangle$ be a repetition-free sequence of agents and let π be any of its permutations:

Weak Commutativity

$$\mathbf{wC}_\ell \quad \Box_{i_1} \dots \Box_{i_\ell} \varphi \rightarrow \Box_{\pi_{i_1}} \dots \Box_{\pi_{i_\ell}} \varphi$$

- We call $\mathbf{C}^b\text{-S5}_n$ the logic S5_n augmented with axiom \mathbf{C}^b .
- We call $\mathbf{wC}_\ell\text{-S5}_n$ the logic S5_n augmented with axiom \mathbf{wC}_ℓ (for all π).

Benefits of Semantic Approach

We obtain positive results:

Logic	Decidability
$K_n, K_n, KT_n, K4_n, K45_n, S4_n, S5_n$	UNDECIDABLE [AB13]
$C^b-S5_n (n>2)$	UNDECIDABLE [Bur+23]
C^b-S5_2	DECIDABLE [Bur+23]
$wC_\ell-S5_n$	
$C-S5_n$	

- Well-known epistemic planning formalism are captured by $C-S5_n$.
- Flexible approach: different axioms can be devised depending on the situation.
- No strong restrictions on modal depth.

Current approaches:

- Compilation of **fragments of DEL** into **classical planning**.
- Bisimulation-contraction techniques.
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Future directions:

- **Symbolic** approaches: SMT encodings, syntactic models.
- Heuristics, heuristics, heuristics.
 - Currently working on: **Epistemic Planning Graph**.
- **Bounded** bisimulation contractions.

Benchmarks for Epistemic Planning

The many fragments of DEL are hard to compare:

- Different ad hoc languages (if any) capture only a part of DEL.
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We need a unified language for the entire DEL semantics. This would allow the following:

- Standard language to represent epistemic planning domains.
- Development of a publicly available and shared set of benchmarks.
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Currently working on **EPDDL**:

- Borrows the well-known syntax of **PDDL** and extends it to capture the whole **DEL semantics**.

Public Announcement

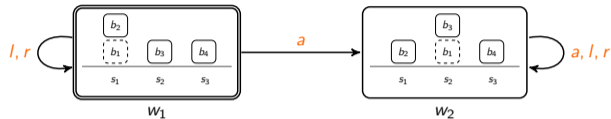
a, l, r



$e : \langle On(b_1, s_3), T \rangle$

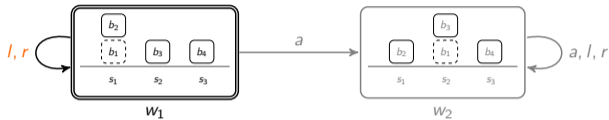
Belief Revision in DEL

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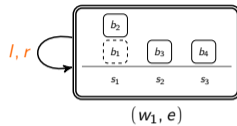


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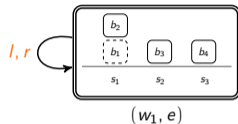
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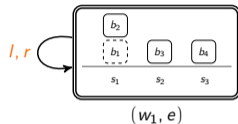
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How do we fix this?

- **Plausibility models**: belief of the agent is captured by the most plausible worlds.
 - We recover **a**'s beliefs by looking at what he considers to be plausible.
- **Recovery**: prior to public announcements, we do a recovering action that “expands” the agents' beliefs.
- **Modifying the product update operator**.

Epistemic planning is still a relatively recent research area.

- Many things still to address.
- Has not been exploited in real scenarios.

Different areas would benefit from epistemic planning and reasoning:

- **Multi-Agent Systems**

- Self-driving vehicles
- Social commitments
- Business Process Management

- **Legal reasoning**

- Virtually any scenario involving uncertainty and/or different perspectives

THANK YOU

Questions?